

WP₁

Security, Trust and Privacy

WP₁ Status

“Security, trust and privacy”

- Timestamping and Blockchain immutability from computational hardness (non-PoW)
- Short cryptographic proofs of predicates
- Cryptographic mechanisms to share data privately
- Identity management, authentication and authorization

- 4y PhD, Aron van Baarsen, 1-3-2020
- (PhD, Esteban Landerreche)

Research output:

- *[LSS20] Non-interactive Cryptographic Timestamping based on Verifiable Delay Functions,* Financial Crypto 2020
- *[LS18] On Immutability of Blockchains,* ERCIM Workshop 2018

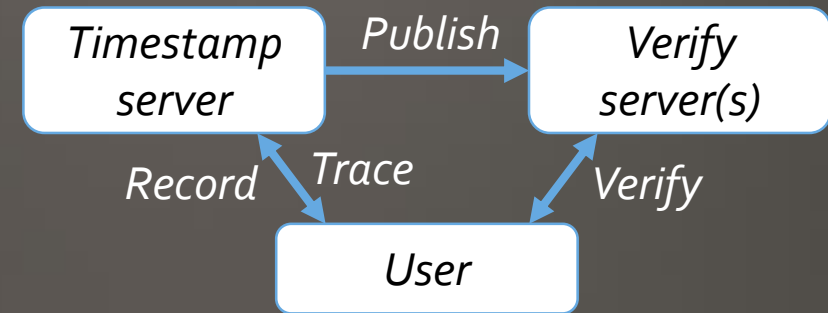
WP₁ Research Output:

Cryptographic Timestamping

Cryptographic Timestamping

Interactive Timestamping:

- Timestamp server + Verify server
- Record a message M at time T_r
 - Output receipt to User
 - Output commitment to Verify server
- Verify timestamp of M at time T_p
 - Timestamp server outputs trace
 - User verifies against Verify server
- Security notions:
 - Backdating security:
Cannot create claim with $T_f < T_r$
 - Postdating security:
Cannot create claim with $T_f > T_r$



- [HS91]
 - Use hash-chains: h_0, h_1, \dots
 - $h_i = \text{Hash}(h_{i-1}, M_i)$
 - Publish all hashes in trusted media
 - Interactive:
Claim must be verified against media
 - Secure if Prover and Verifier don't collude
- [BHS93]:
 - Batch messages in blocks using Merkle Tree

Blockchain Immutability & Timestamping

- Distributed Ledgers (DLT) should be immutable
 - Sufficiently old Transactions cannot be removed, altered, or inserted
 - [Nako8] Bitcoin construction:
 - Hash-chains: any change in a block M_i affects all subsequent hashes
 - Proof-of-Work: for each block solve *parallel computational problem*
 - [Nako8] Immutable against attackers with <50% computational power
- [GMG15]: Timestamping service build on Bitcoin
- [LS18,LSS20]: (Non-interactive) timestamping based on *serial computational problem*

Medium-hard Computational Problems for Timestamping

- Proof-of-Work

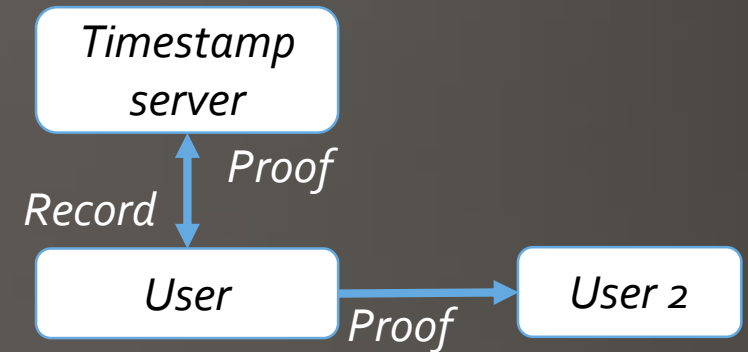
- Input block M , difficulty D
- Find N such that $\text{Hash}(M) < D$
- Massively parallel
- Bitcoin: 2^{67} H/sec, 2^{91} H/year
- Massively wasteful: ~ 65 TWh/year
- Comparable to power use of Czech Republic
- Can be used for Bitcoin consensus

- VDF: Verifiable Delay Function

- Input block M , iterations T
- Computes T serial operations in T wallclock time
- Generates short efficient proof of correctness
- Inherently sequential
- Attacks require much higher speed than honest parties
- Physical limitations on iterations/sec
- Need optimized VDF cores for honest usage
- Can't be used for Bitcoin consensus

(Non-interactive) timestamping using VDFs

- Non-interactive Timestamping:
 - Only need Trusted Timestamping server (no Verify server needed)
 - Record a message M at time T_r
 - Output receipt to User
 - Verify timestamp of M at time T_p
 - Timestamp server outputs non-interactive proof of age ($T_p - T_r$)
 - User verifies proof
 - Proof is over relative time, not absolute time
 - Absolute time can be achieved through online timestamping service

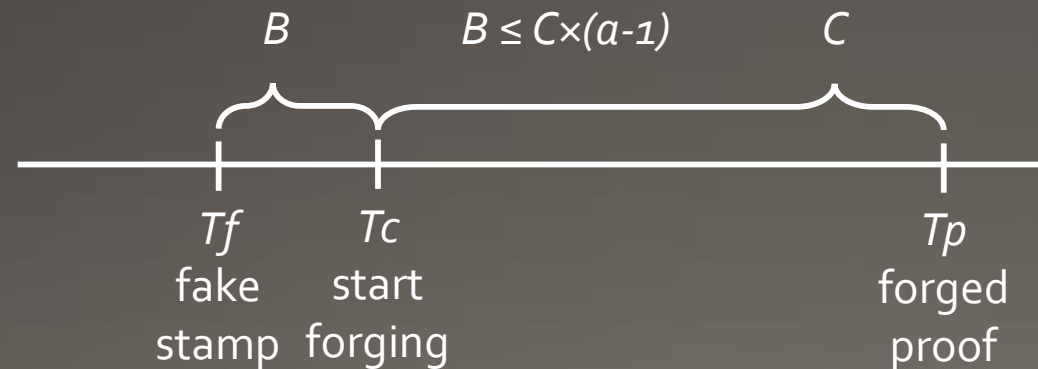


(Non-interactive) timestamping using VDFs

- [LS18,LSS20] Trusted Timestamping Service:
 - Hash-chain h_0, h_1, \dots
 - $h_i = \text{Hash}(h_{i-1}, M_i, \text{time}, \text{VDF}(h_{i-1}, T_i), \text{sig})$
 - Continuously compute VDF from last message until next message
- Record a message M at time T_r :
 - Stop ongoing VDF and finalize VDF proof
 - Compute signature and next hash in chain
 - Continue VDF from new hash
- Non-interactive timestamping proof for M :
 - Output hash-chain from T_r onwards
 - Copy ongoing VDF state, finalize and output VDF proof

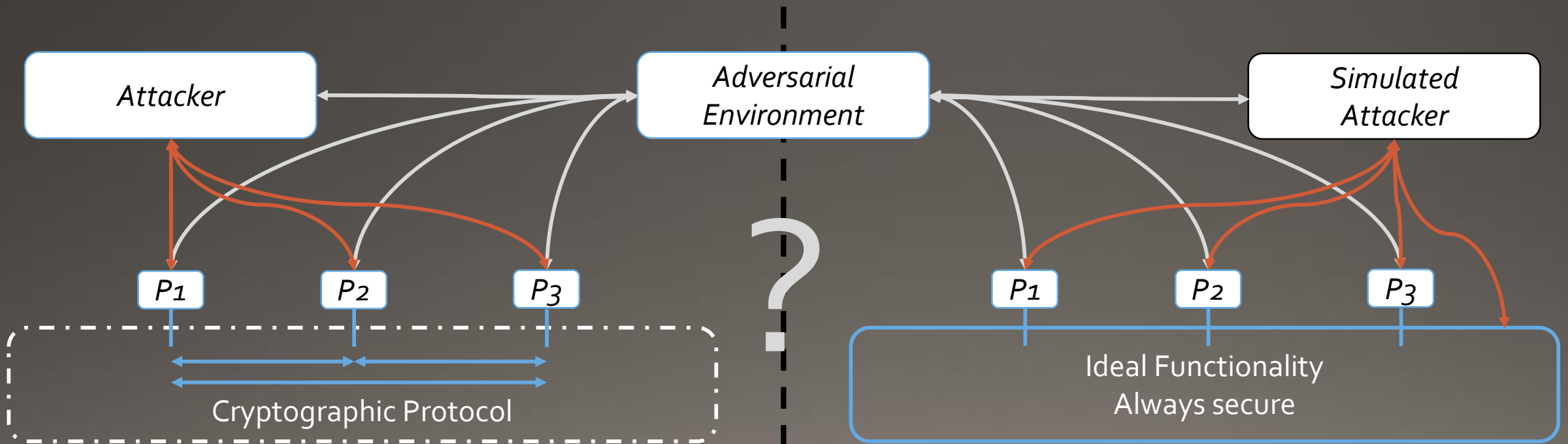
(Non-interactive) timestamping using VDFs

- [LSS20] Proven security bounding backdating attacks
 - Requires control signing key
 - Requires faster VDF core by factor $a > 1$
 - E.g. $a = 1.1$ then after 10 days at most backdated $10 \times (a - 1) = 1$ day



(Non-interactive) timestamping using VDFs

- Proven secure in *Universal Composability Framework*
 - A Gold standard in Provable-Secure / Secure-by-Design protocols
 - Protocol is secure if no adversarial environment can distinguish between the real world (left) and the ideal world (right)



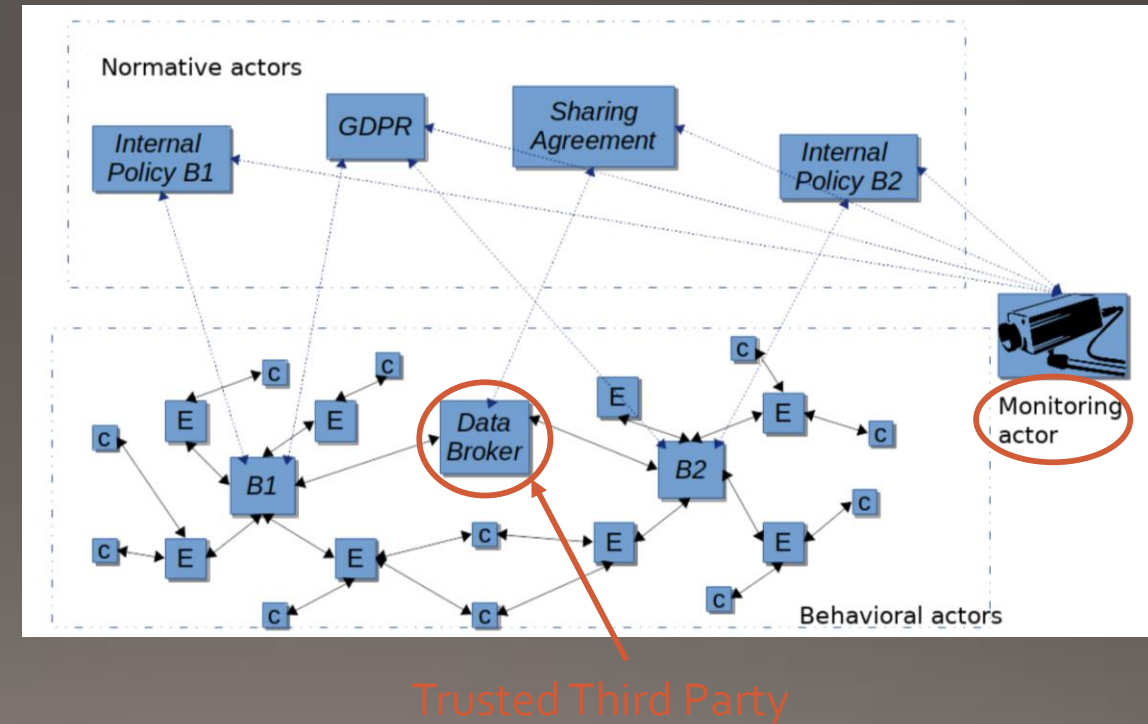
(Non-interactive) timestamping using VDFs

- Based on VDF => Only 1 VDF core sufficient
- Non-interactive proof based on hash-chains
 - => Can be integrated into any other hash-chain protocol, such as DLT protocols at the block level
 - => DLT immutability based on computational problem
 - Changing one block in DLT chain requires recomputing all subsequent VDF proofs
- [LS18] *On Immutability of Blockchains*
 - Shows protocol combining N timestamping servers
 - Adds threshold security: attackers needs to corrupt majority of N servers

WP₁ Plans

WP₁ Plans

- Activity 3 (ongoing)
 - Follow-up work on Timestamping
 - Prove secure VDF construction in Universal Composability Framework
 - Stronger foundations for underlying mathematical groups with hidden order
- Future activity 4
 - Extend eFLINT w/ *Trust Management*
 - Integration with WP₃
- Future activity 5
 - Trusted third party “Data broker”
 - Replace with cryptographic protocol: *Secure Multi-Party Computation*
 - Integration with WP₂



WP₁ Plans

Activity 4: Identity management, authentication and authorization

- Goal: eFLINT extension Trust Management
 - Inspiration: KeyNote, Blaze et al., '98
- DSL Extensions:
 - Principals (entities)
 - Credentials (pubkey, authorization, delegation)
 - Trust policy
 - Internal crypto processing
- Integration with WP₃

WP₁ Plans

Activity 5: Cryptographic mechanisms to share data privately

- eFLINT Demo uses Trusted Party for risk computation on multi-party data
- Goal: replace by *Secure Multi-Party Computation*
 - Generically possible but generally inefficient
 - Construct optimized protocol suitable for KYC usecase
- Strong guarantees:
 - Only output is learned
 - No other information is leaked
 - Unless $> 1/3$ of parties collude
- Integration with WP₂: High-Performance Accelerators